Most real networks are dynamic

- Facebook friendship
 - People joining/leaving
 - Friend/Unfriend
- Twitter mention network
 - Each mention has a timestamp
 - If aggregated: every day/month/year, #messages change
- Soccer pass network (during match, between match)
- Internet infrastructure
- <u>۰</u>...

- Most real world networks evolve
 - Nodes can appear/disappear
 - Edges can appear/disappear
 - Nature of relations can change
- How to represent those changes?
- How to manipulate dynamic networks?

Semantic level

Relations

Long term

-Friend -Colleague -Family relation

- - - -

Short term ?

-Collaborators in the same project -Same team in a game -Attendees of the same meeting

- . . .

Interactions

Instantaneous

-e-mail -Text message -Co-authoring

. . .

With duration

-Phone call -Discussion in real life -Participate in a same meeting











- Exemple in practice: Sociopattern dataset
 - ► Every 20s, list of individuals at distance ≈ 1,5m
 - Dataset : sequence of graphs or temporal edge list

353304 00	48	644
353304 00	6 3	672
353304 00	656	682
353304 00	632	67
353304 20	1492	1613
353304 20	656	682
353304 20	1632	1671
1353304140	1148	1644
353304 60	656	682
353304 60	1108	1601
353304 60	1632	1671
353304 60	626	698











DYNAMIC NETWORKS To sum up

- Collected dataset of relations (Facebook "friends", affiliations, etc.)
 - Can be discretized for analysis if needed
- Collected dataset of interactions (e-mail, face to face, etc)
 - Can be aggregated in series of graphs
 - Can be transformed in intervals (relations or lasting interactions)
 - Can be studied as it is

ANALYZING DYNAMIC NETWORKS

UNSTABLE SNAPSHOTS

UNSTABLE SNAPSHOTS

- The evolution is represented as a series of a few snapshots.
- Many changes between snapshots
 - Cannot be plotted as a "movie"



UNSTABLE SNAPSHOTS

- Each snapshot can be studied as a static graph
- The evolution of the properties can be studied "manually"
- "Node X had low centrality in snapshot t and high centrality in snapshot t+n"

Edges correspond to relations

- They change slowly
- The network is well defined at any t
 - Temporal network: nodes/edges described by (long lasting) intervals
 - Enough snapshots to track nodes
- A static analysis at every (relevant) t gives a dynamic vision
- No formal distinction with previous case

- Visualization
 - Problem of stability of node positions





[Leskovec: Graph Evolution: Densification and Shrinking Diameters]

Centralities



TIME SERIES ANALYSIS

- TS analysis is a large field of research
- Time series: evolution of a value over time
 - Very popular for stock market analysis...
- Killer applications:
 - Detection of periodic patterns
 - Detection of anomalies
 - Identification of global trends
 - Evaluation of auto-correlation
 - Prediction of future values
- e.g. ARIMA (Autoregressive integrated moving average)
 - From statsmodels.tsa.arima_model import ARIMA

- Temporal Network of interactions
 - Link stream (twitter messages,...)
 - Interval graph (Physical interactions, ...)
- The network at a given t is not meaningful
- How to analyze such a network?





- Until recently, network was transformed using snapshot/sliding windows
- Tools developed to deal with such networks

- [Holme 2012]: mostly about paths, walks, distances...
- [Latapy 2018]: Other things (centralities, ...)
- Idea: Generalize all graphs definitions to temporal networks
- => If all nodes and all edges always present, same values as for a static graph

CENTRALITIES & NETWORK PROPERTIES IN STREAM GRAPHS

stream graph
$$S = (T, V, W, E)$$

T: Possible Time V: vertices W: Vertices presence time E: Edges presence time

Number of nodes:

Total presence of nodes

Total dataset duration

(not an integer value...)

$$n = \sum_{v \in V} n_v = \frac{|W|}{|T|}$$

e.g.: 2 if 4 nodes half the time

Number of edges:

Total presence of edges

Total dataset duration

(not an integer value...)

 $m = \sum_{uv \in V \otimes V} m_{uv} = \frac{|E|}{|T|}$

e.g.: I if I edge all the time

Neighborhood of a node $N(v) = \{(t, u), (t, uv) \in E\}$

$$d(v) = \frac{|N(v)|}{|T|} = \sum_{u \in V} \frac{|T_{uv}|}{|T|}$$



Figure 5: Two examples of neighborhoods and degrees of nodes. We display in black the links involving the node under concern, and in grey the other links. Left: $N(a) = ([1,3] \cup [7,8]) \times \{b\} \cup [4.5,7.5] \times \{c\}$ is in blue, leading to $d(a) = \frac{3}{10} + \frac{3}{10} = 0.6$. Right: $N(c) = [2,5] \times \{a\} \cup [1,8] \times \{b\} \cup [6,9] \times \{d\}$ is in blue, leading to $d(c) = \frac{13}{10} = 1.3$.

Average node degree

$$d(V) = \frac{1}{n} \cdot \sum_{v \in V} n_v \cdot d(v) = \sum_{v \in V} \frac{|T_v|}{|W|} \cdot d(v)$$

Clustering coefficient of a node

$$cc(v) = \delta(N(v)) = \frac{\sum_{uw \in V \otimes V} |T_{vu} \cap T_{vw} \cap T_{uw}|}{\sum_{uw \in V \otimes V} |T_{vu} \cap T_{vw}|}$$

Probability that if we take 2 random neighbors at random time, they are linked

$$\delta(S) = \frac{\sum_{uv \in V \otimes V} |T_{uv}|}{\sum_{uv \in V \otimes V} |T_u \cap T_v|} = \frac{\int_{t \in T} |E_t| \, \mathrm{d}t}{\int_{t \in T} |V_t \otimes V_t| \, \mathrm{d}t}$$

Density: probability if we take a random pair of nodes at random time that there is an edge

$$\delta(S) = \frac{\sum_{uv \in V \otimes V} |T_{uv}|}{\sum_{uv \in V \otimes V} |T_u \cap T_v|} = \frac{\int_{t \in T} |E_t| \, \mathrm{d}t}{\int_{t \in T} |V_t \otimes V_t| \, \mathrm{d}t}$$

Total edge presence

e.g.: 10 if 2 edges present over 5 periods



Total **overlapping time** between each pair of nodes =>An edge is possible



Figure 2: Two stream graphs with n = 2 nodes, m = 1 link, but with different densities: Left: $\delta = 0.75$. Right: $\delta = 1$.

- Note that we can define particular cases of density:
 - Density of a pair of nodes
 - Density of a node

$$\delta(uv) = \frac{|T_{uv}|}{|T_u \cap T_v|}, \quad \delta(v) = \frac{\sum_{u \in V, u \neq v} |T_{uv}|}{\sum_{u \in V, u \neq v} |T_u \cap T_v|}$$

A clique of graph G is a cluster C of G of density 1. In other words, all pairs of nodes involved in C are linked together in G. A clique C is maximal if there is no other clique C' such that $C \subset C'$.



PATHS AND DISTANCES IN STREAM GRAPHS

PATHS

- A path in a stream graphs
 - Starts at a node and a date
 - Ends at a node and a date
 - Has a length (number of hops)
 - Has a duration (duration from leaving node to reaching node)



- Several types of shortest paths in Stream graphs:
 - Shortest path: minimal length
 - Fastest path: minimal duration
 - Foremost path: first to reach
 - Fastest shortest paths
 - Minimum duration among minimal length
 - Shortest fastest paths
 - Minimal length among minimal duration



Shortest paths from (I, d) to (9, c) ?



Shortest paths from (1, d) to (9, c) ? =>e.g. (2.5,d,b)(3,b,a)(7,a,c)





(3,d,b),(3,b,a),(4.5,a,c)





...(4.5,a,c)



Fastest shortest path from (I, d) to (9, c)?



Fastest shortest path from (I, d) to (9, c)?

(3, d, b), (3, b, a), (4.5, a, c)



Shortest Fastest path from (I, d) to (9, c) ?

OTHER DEFINITIONS ON STREAM GRAPHS

CONNECTED COMPONENTS

- Weakly connected component:
 - There is at least a non-temporally respecting path



CLOSENESS - BETWEENESS

$$\mathcal{C}_t(v) = \sum_{u \in V} \int_{\substack{s \in T \\ (s,u) \neq (t,v)}} \frac{1}{c_t(v, (s, u))} \, \mathrm{d}s$$

$$\mathcal{B}(t,v) = \sum_{u \in V, w \in V} \int_{i \in T_u, j \in T_w} \frac{\sigma((i,u), (j,w), (t,v))}{\sigma((i,u), (j,w))} \,\mathrm{d}i \,\mathrm{d}j$$

RANDOM MODELS FOR DYNAMIC NETWORKS

RANDOM MODELS

- In many cases, in network analysis, useful to compare a network to a randomized version of it
 - Clustering coefficient, assortativity, modularity, ...
- In a static graph, 2 main choices:
 - Keep only the number of edges (ER model)
 - Keep the number of edges and the degree of nodes (Configuration model)
- In dynamic networks, it is more complex...

RANDOM MODELS

- [Gauvin 2018]
- Four families of shuffling:
 - Link Shuffling
 - => Conserve time of events by nodes and edges, switch end points
 - Timeline shuffling
 - => Switch time of events between elements (e.g. activations of nodes a and b)
 - Snapshot shuffling
 - =>Keep the order of snapshots, randomize network inside snapshot
 - Sequence Shuffling
 - =>Keep each snapshot identical, but switch randomly their order
- Shufflings can be combined...

RANDOM MODELS



TO SUM UP

CONCLUSION

- Currently, most people still use the snapshot approaches
- Most researchers agree that it has many drawbacks
- But currently:
 - No single framework
 - No library
 - Dataset not has ubiquitous as static graphs
 - (often privates...)

PRACTICALS

PRACTICALS

- I)Select a temporal networks composed of snapshots
 2)Plot the evolution of centralities
- 3) Try to smooth out your network by merging snapshots
 4) Recompute evolution centralities
- 5)Compare this evolution with a snapshot randomized network (Randomize networks, keep snapshot orders)
 - 6)Is the tendency correlated with tendencies in the random network?
- 7)Compute the edge density of pairs of nodes, for nodes that are not present all the time. What do you observe? Does it makes sense? In which case is it important?

PRACTICALS

• Datasets:

- On my webpage, Game of thrones, 2 possibilities:
 - Books: One network by book. Stable. Can study differences
 - Series: One SN every 10 scenes. Unstable. More interesting (in particular for time series analysis), more difficult
- There Series:
- https://figshare.com/articles/TV_Series_Networks_of_characters/2199646/11